

A polynomial algorithm to check the fixed point property for ordered sets of dimension 2

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Definition.

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n Open Questions

Definition. An **ordered set** is a pair (P, \leq) of a set P and a reflexive, antisymmetric and transitive relation \leq , the order relation.

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Fixed Point Property

d2-collapsibility

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Definition. A finite ordered set P is **two-dimensional** iff its order is the intersection of two linear orders.



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Definition. A function from a finite ordered set (P, \leq) to another ordered set (Q, \leq) is called **order-preserving** *iff*, for all $x, y \in P$ we have that $x \leq y$ implies $f(x) \leq f(y)$.

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Definition. A finite ordered set is said to have the **fixed point property** iff every order-preserving self-map has a fixed point.

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Definition.

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Definition. *Let P be an ordered set.*

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Definition. Let P be an ordered set. Then an order-preserving map $r : P \rightarrow P$ is called a **retraction** iff $r^2 = r$

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Theorem. *Let P be an ordered set with the fixed point property and let* $r : P \rightarrow P$ *be a retraction.*

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Theorem. Let *P* be an ordered set with the fixed point property and let $r : P \to P$ be a retraction. Then r[P] has the fixed point property.

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What about possible converses?

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Definition. Let *P* be a finite ordered set and let $x, y \in P$. If x < y and there is no $z \in P$ so that x < z < y, then y is called an **upper cover** of x

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Definition. Let *P* be a finite ordered set and let $x, y \in P$. If x < y and there is no $z \in P$ so that x < z < y, then *y* is called an **upper cover** of *x* and *x* is called a **lower cover** of *y*.

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Definition. Let P be a finite ordered set and let $x, y \in P$. If x < y and there is no $z \in P$ so that x < z < y, then y is called an **upper cover** of x and x is called a lower cover of y. In a finite ordered set, a point is called irreducible iff it has exactly one upper cover or exactly one lower cover.

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Theorem (Rival, 1976).

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Theorem (Rival, 1976). *Let* P *be a finite ordered set and let* $a \in P$ *be irreducible.*

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Theorem (Rival, 1976). *Let* P *be a finite ordered set and let* $a \in P$ *be irreducible. Then* P *has the fixed point property iff*

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Theorem (Rival, 1976). Let P be a finite ordered set and let $a \in P$ be irreducible. Then P has the fixed point property iff $P \setminus \{a\}$ has the fixed point property.

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Definition.

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Definition. A finite ordered set P is called dismantlable iff

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Definition. A finite ordered set P is called **dismantlable** iff |P| = 1 or there is a point $x \in P$ such that 1. x is irreducible in P, 2. $P \setminus \{x\}$ is dismantlable.

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Definition.

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Definition. Let *P* be a finite ordered set.

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Definition. Let P be a finite ordered set. A point $a \in P$ is called **retractable** (to the **point** $b \in P \setminus \{a\}$)

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Definition. Let *P* be a finite ordered set. A point $a \in P$ is called **retractable (to the point** $b \in P \setminus \{a\}$) iff for all $x \in P, x > a$ implies $x \ge b$

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d2-collapsibility

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Fixed Point Property Retractable Points Dimension 2 d2-collapsibility The Algorithm Open Questions

Theorem (BS, 1993).

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Fixed Point Property

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Theorem (BS, 1993). *Let* P *be a finite ordered set and let* $a \in P$ *be retractable.*

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Theorem (BS, 1993). *Let* P *be a finite ordered set and let* $a \in P$ *be retractable. Then* P *has the fixed point property iff*

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Theorem (BS, 1993). Let P be a finite ordered set and let $a \in P$ be retractable. Then P has the fixed point property iff 1. $P \setminus \{a\}$ has the fixed point property.

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Theorem (BS, 1993). Let P be a finite ordered set and let $a \in P$ be retractable. Then P has the fixed point property iff

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Definition.

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Theorem (BS, 1993). Let P be a finite ordered set and let $a \in P$ be retractable. Then P has the fixed point property iff 1. $P \setminus \{a\}$ has the fixed point property. 2. $\uparrow a \setminus \{a\}$ has the fixed point property.

Definition. A finite ordered set P is called **connectedly collapsible** *iff*

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Definition. A finite ordered set P is called **connectedly** collapsible iff |P| = 1

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Definition. A finite ordered set P is called **connectedly** collapsible iff |P| = 1 or there is a point $x \in P$ such that

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d2-collapsibility

An Example (Rutkowski)

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An Example (Rutkowski)

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Dimension 2

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Theorem (Fofanova, Rutkowski, Rival, 1994).

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Theorem (Fofanova, Rutkowski, Rival, 1994). Let P be a finite ordered set of interval dimension 2.

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Theorem (Fofanova, Rutkowski, Rival, 1994). Let P be a

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Theorem (Fofanova, Rutkowski, Rival, 1994). Let P be a finite ordered set of interval dimension 2. Then P contains a point with a unique lower cover that is minimal in P, or, P contains a minimal element a that is retractable to another minimal element b.

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Definition.

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Definition. *Let P be a finite ordered set and let* $a \in P$ *. Then a is called* **d2-retractable** (*in P*) *iff one of the following holds.*

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Definition. *Let* P *be a finite ordered set and let* $a \in P$ *. Then a is called* **d2-retractable** (*in* P) *iff one of the following holds.*

1. The point a is irreducible, or,

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Theorem (Fofanova, Rutkowski, Rival, 1994). Let P be a finite ordered set of interval dimension 2. Then P contains a point with a unique lower cover that is minimal in P, or, P contains a minimal element a that is retractable to another minimal element b.

Definition. Let P be a finite ordered set and let $a \in P$. Then a is called **d2-retractable** (in P) iff one of the following holds.

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Definition. *Let P be a finite ordered set. Then P is said to have a* **d2-retraction sequence**

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- 1. The point a is irreducible, or,
- 2. The point a is not irreducible, but minimal in P and retractable to another minimal element b of P.

Definition. Let P be a finite ordered set. Then P is said to have a **d2-retraction sequence** iff there is a retraction sequence $a_1, \ldots, a_{|P|}$ so that, for all $i \in \{1, \ldots, |P| - 1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

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Lemma (BS, henceforth omitted, joke gets old).

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Lemma (BS, henceforth omitted, joke gets old). *Let P be a finite ordered set that has a d2-retraction sequence and let* $r : P \rightarrow P$ *be a retraction.*

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Lemma (BS, henceforth omitted, joke gets old). *Let P be a finite ordered set that has a d2-retraction sequence and let* $r : P \rightarrow P$ *be a retraction. Then* r[P] *has a d2-retraction sequence.*

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Lemma.

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Lemma (BS, henceforth omitted, joke gets old). *Let P be a finite ordered set that has a d2-retraction sequence and let* $r : P \rightarrow P$ *be a retraction. Then* r[P] *has a d2-retraction sequence.*

Lemma. *Let P be a finite ordered set and let* $x \in P$ *be d2-retractable.*

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Lemma (BS, henceforth omitted, joke gets old). *Let P be a finite ordered set that has a d2-retraction sequence and let* $r : P \rightarrow P$ *be a retraction. Then* r[P] *has a d2-retraction sequence.*

Lemma. Let P be a finite ordered set and let $x \in P$ be d2-retractable. Then P has a d2-retraction sequence iff $P \setminus \{x\}$ does.

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Lemma.

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Lemma. Let P be a finite ordered set and let $x \in P$ be d2-retractable. Then P has a d2-retraction sequence iff $P \setminus \{x\}$ does.

Lemma. For any ordered set P, it takes at most $3|P|^3$ steps to decide if there is a d2-retraction sequence for P

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Lemma (BS, henceforth omitted, joke gets old). *Let P be a finite ordered set that has a d2-retraction sequence and let* $r : P \rightarrow P$ *be a retraction. Then* r[P] *has a d2-retraction sequence.*

Lemma. Let P be a finite ordered set and let $x \in P$ be d2-retractable. Then P has a d2-retraction sequence iff $P \setminus \{x\}$ does.

Lemma. For any ordered set P, it takes at most $3|P|^3$ steps to decide if there is a d2-retraction sequence for P and, if so, to compute a d2-retraction sequence for P.

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Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Lemma.

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Lemma. Let P be a finite ordered set that has a d2-retraction sequence $a_1, \ldots, a_{|P|}$.

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Lemma. Let P be a finite ordered set that has a d2-retraction sequence $a_1, \ldots, a_{|P|}$. Then P has the fixed point property iff

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Lemma. Let P be a finite ordered set that has a d2-retraction sequence $a_1, \ldots, a_{|P|}$. Then P has the fixed point property iff, for all $i \in \{1, \ldots, |P| - 1\}$ so that, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, a_i is minimal, not irreducible, but retractable to minimal element of $P \setminus \{a_1, \ldots, a_{i-1}\}$, we have that $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i$ has the fixed point property.

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Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Lemma.

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Open Questions

Lemma. *Let P be a finite ordered set.*

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Lemma. Let P be a finite ordered set. Then P is called **d2-collapsible** iff P is empty, a singleton

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Lemma. Let P be a finite ordered set. Then P is called **d2-collapsible** iff P is empty, a singleton, or, |P| > 1 and P has a d2-retraction sequence $a_1, \ldots, a_{|P|}$ so that

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Lemma. Let P be a finite ordered set. Then P is called **d2-collapsible** iff P is empty, a singleton, or, |P| > 1 and P has a d2-retraction sequence $a_1, \ldots, a_{|P|}$ so that, for each $i \in \{1, \ldots, |P| - 1\}$, if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element of $P \setminus \{a_1, \ldots, a_{i-1}\}$

d2-collapsible: *Don't allow the empty set.*

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Proposition.

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Proposition. *Let P be a finite ordered set.*

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Proposition. *Let P be a finite ordered set.*

1. If P is d2-connectedly collapsible, then P is d2-collapsible.

Proposition. *Let P be a finite ordered set.*

- 1. If P is d2-connectedly collapsible, then P is d2-collapsible.
- 2. *If P is d2-connectedly collapsible, then P has the fixed point property.*

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Proposition. Let P be a finite ordered set.

- 1. If P is d2-connectedly collapsible, then P is d2-collapsible.
- 2. *If P is d2-connectedly collapsible, then P has the fixed point property.*
- 3. *If P is d2-collapsible, then P has the fixed point property iff P is d2-connectedly collapsible.*

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Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Lemma.

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Lemma. Let *P* be a d2-collapsible ordered set and let $r : P \rightarrow P$ be a retraction.

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Lemma. Let *P* be a d2-collapsible ordered set and let $r : P \to P$ be a retraction. Then r[P] is d2-collapsible.

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Lemma. Let P be a d2-collapsible ordered set and let $r : P \rightarrow P$ be a retraction. Then r[P] is d2-collapsible. Moreover, if P is d2-connectedly collapsible, then r[P] is d2-connectedly collapsible.

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Lemma. Let P be a d2-collapsible ordered set and let $r : P \to P$ be a retraction. Then r[P] is d2-collapsible. Moreover, if P is d2-connectedly collapsible, then r[P] is d2-connectedly collapsible.

Lemma.

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P.

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element of $P \setminus \{a_1, \ldots, a_{i-1}\}$

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible.

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, ..., a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, ..., |P| - 1\}$, we have that if, in $P \setminus \{a_1, ..., a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element of $P \setminus \{a_1, ..., a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, ..., a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, ..., a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, ..., |P| - 1\}$, we have that if, in $P \setminus \{a_1, ..., a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element of $P \setminus \{a_1, ..., a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, ..., a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."

Problem: Recursively, that's a lot of sets.

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Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Definition.

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Definition. *Let C be a finite ordered set.*

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Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

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Open Questions

Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

Theorem.

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Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

Theorem. (*Duffus-Poguntke-Rival 1980*, *Farley 1993*) Every finite ordered set has a unique core.

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Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

Theorem. (*Duffus-Poguntke-Rival 1980, Farley 1993*) Every finite ordered set has a unique core.

Lemma.

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Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

Theorem. (*Duffus-Poguntke-Rival 1980, Farley 1993*) Every finite ordered set has a unique core.

Lemma. *Let P be a finite ordered set. Then P is d2-collapsible iff its core is d2-collapsible.*

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Definition. Let C be a finite ordered set. Then C is called a **core** iff P contains no irreducible elements.

Theorem. (*Duffus-Poguntke-Rival 1980, Farley 1993*) Every finite ordered set has a unique core.

Lemma. Let P be a finite ordered set. Then P is d2-collapsible iff its core is d2-collapsible. The same result holds for d2-connectedly collapsible ordered sets.

Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Lemma.

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Lemma. Let *H* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $a \in \uparrow x \cap \uparrow y$ be irreducible in *P*.

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Fixed Point Property

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Lemma. Let *H* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $a \in \uparrow x \cap \uparrow y$ be irreducible in *P*. Then *a* is irreducible in $\uparrow x \cap \uparrow y$.

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Lemma. Let *H* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $a \in \uparrow x \cap \uparrow y$ be irreducible in *P*. Then *a* is irreducible in $\uparrow x \cap \uparrow y$.

Definition.

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Lemma. Let *H* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $a \in \uparrow x \cap \uparrow y$ be irreducible in *P*. Then *a* is irreducible in $\uparrow x \cap \uparrow y$.

Definition. *Let P be a finite ordered and let* $Q \subseteq P$ *.*

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Lemma. Let *H* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $a \in \uparrow x \cap \uparrow y$ be irreducible in *P*. Then *a* is irreducible in $\uparrow x \cap \uparrow y$.

Definition. *Let P be a finite ordered and let* $Q \subseteq P$ *. Then P has a* **retraction sequence to** *Q iff*

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Definition. Let P be a finite ordered and let $Q \subseteq P$. Then P has a **retraction sequence to** Q iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

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Definition. Let *P* be a finite ordered and let $Q \subseteq P$. Then *P* has a **retraction sequence to** *Q* iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

Lemma.

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Definition. Let P be a finite ordered and let $Q \subseteq P$. Then P has a **retraction sequence to** Q iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

Lemma. Let P be a finite ordered set, let $x, y \in P$ be two incomparable elements

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Definition. Let P be a finite ordered and let $Q \subseteq P$. Then P has a **retraction sequence to** Q iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

Lemma. Let *P* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $Q \subseteq P$ be a subset of *P* that contains *x* and *y*

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Definition. Let P be a finite ordered and let $Q \subseteq P$. Then P has a **retraction sequence to** Q iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

Lemma. Let P be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $Q \subseteq P$ be a subset of P that contains x and y and for which P has a d2-retraction sequence to Q.

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Definition. Let P be a finite ordered and let $Q \subseteq P$. Then P has a **retraction sequence to** Q iff the elements of $P \setminus Q$ can be arranged in a sequence $a_1, \ldots, a_{|P|-|Q|}$ so that, for all $i \in \{1, \ldots, |P|-1\}$, the point a_i is d2-retractable in $P \setminus \{a_1, \ldots, a_{i-1}\}$.

Lemma. Let *P* be a finite ordered set, let $x, y \in P$ be two incomparable elements and let $Q \subseteq P$ be a subset of *P* that contains *x* and *y* and for which *P* has a d2-retraction sequence to *Q*. Then the core of $\uparrow_Q x \cap \uparrow_Q y$ is isomorphic to the core of $\uparrow x \cap \uparrow y$.

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Fixed Point Property	Retractable Points	Dimension 2	The Algorithm	Open Questions

Lemma.

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P.

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Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff

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ity The Algorithm

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P|-1\}$, we have that

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Open Questions

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$

Dimension 2

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Open Questions

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible.

Dimension 2

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Open Questions

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."

Dimension 2

d2-collapsibil

ity The Algorithm

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i \setminus \{a_i\}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_i \cap \uparrow_{P \setminus \{a_1, \ldots, a_{i-1}\}} a_{i+j}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."

Dimension 2

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ility The Algorithm

Lemma. Let P be a finite ordered set with a d2-retraction sequence and let $a_1, \ldots, a_{|P|}$ be any d2-retraction sequence for P. Then P is d2-collapsible iff, for any $i \in \{1, \ldots, |P| - 1\}$, we have that if, in $P \setminus \{a_1, \ldots, a_{i-1}\}$, the point a_i is minimal, not irreducible, but retractable to another minimal element a_{i+j} of $P \setminus \{a_1, \ldots, a_{i-1}\}$, then $\uparrow a_i \cap \uparrow a_{i+j}$ is d2-collapsible. The same result holds when "d2-collapsible" is replaced with "d2-connectedly collapsible."



Algorithm.

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The Algori

Open Questions

Algorithm. The algorithm D2CCcheck.

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1. Input: The algorithm takes as input a nonempty list (P_1, \ldots, P_n) of ordered sets.

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 Input: The algorithm takes as input a nonempty list (P₁,...,P_n) of ordered sets.
Initialization.

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- 1. Input: The algorithm takes as input a nonempty list (P_1, \ldots, P_n) of ordered sets.
- 2. Initialization.
 - 2.1 Initialize the list \mathscr{S} to be the input, $\mathscr{S} := (P_1, \ldots, P_n)$.

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- 1. Input: The algorithm takes as input a nonempty list (P_1, \ldots, P_n) of ordered sets.
- 2. Initialization.
 - 2.1 Initialize the list \mathscr{S} to be the input, $\mathscr{S} := (P_1, \ldots, P_n)$.
 - 2.2 Initialize the list \mathcal{D} to be the empty list.

- 1. Input: The algorithm takes as input a nonempty list (P_1, \ldots, P_n) of ordered sets.
- 2. Initialization.
 - 2.1 Initialize the list \mathscr{S} to be the input, $\mathscr{S} := (P_1, \ldots, P_n)$.
 - 2.2 Initialize the list \mathcal{D} to be the empty list.
 - 2.3 Set the counter c := 1.

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3. If $\mathscr{S} = (Q_1, \ldots, Q_m)$ and c > m, return \mathscr{S} , \mathscr{D} , exit.

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If \$\mathcal{S}\$ = (Q₁,...,Q_m) and c > m, return \$\mathcal{S}\$, \$\varnothing\$, exit.
Given \$\mathcal{S}\$ = (Q₁,...,Q_m), consider the set Q_c.

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- 3. If $\mathscr{S} = (Q_1, \dots, Q_m)$ and c > m, return \mathscr{S} , \mathscr{D} , exit.
- 4. Given $\mathscr{S} = (Q_1, \dots, Q_m)$, consider the set Q_c . 4.1 If Q_c is a singleton, append $d_c := TRUE$ to \mathscr{D} .

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3. If $\mathscr{S} = (Q_1, \ldots, Q_m)$ and c > m, return \mathscr{S} , \mathscr{D} , exit.

4. Given $\mathscr{S} = (Q_1, \dots, Q_m)$, consider the set Q_c . 4.1 If Q_c is a singleton, append $d_c := TRUE$ to \mathscr{D} .

4.2 If Q_c is an antichain and $|Q_c| > 1$, append $d_c := AC$ to \mathcal{D} .

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A polynomial algorithm to check the fixed point property for ordered sets of dimension 2

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- 5. Increase c by 1 and continue at 3.

Fixed Point Property	Retractable Points	Dimension 2	d2-collapsibility	Open Questions

Theorem.

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d2-collapsibility

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- 3. *P* is not d2-collapsible iff $\mathcal{D}(P)$ has an entry NO.

d2-collapsibility

The Algorithm

Final Comments

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- 2. Let *P* be a finite ordered set that does not contain any crowns with more than 4 elements. Must *P* be d2-collapsible?